

Uncertainties in Modeling & Simulation

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Learning Objectives

□To understand fundamental concepts of Modeling & Simulation (M&S)

To understand the major sources of epistemic uncertainty in M&S



How to study a system





What is Modeling?



Why mathematical modeling?

□Advantages

Disadvantages



An example of modeling

□Free fall model





Mathematical model

□*Dependent variable=f*(*Independent variable*) y = f(x)

□High dimensional

$$y = f\left(x_1, x_2, \ldots\right)$$

□Parametric systems

$$y = f(x, u)$$

"Noisy" systems

$$y = f(x, u, \gamma)$$



Complexity of Mathematical Models





Model Taxonomy





Modeling & Simulation at Multiple Scales



Various methods used to simulate at different length and time scales



Simulation-based Design



Two types of "uncertainties"

□Aleatory uncertainty (**variability**, irreducible uncertainty, random error)

- inherently associated with the randomness/fluctuation (e.g. environmental stochasticity, inhomogeneity of materials, fluctuation of measuring instruments)
- can only be reduced by taking average of multiple measurements.
- Epistemic uncertainty (incertitude, reducible uncertainty, systematic error)
 - imprecision comes from scientific ignorance, inobservability, lack of knowledge, etc.
 - can be reduced by additional empirical effort (such as calibration).



Random Error

Determines the *precision* of any measurement
 Always present in every physical

- measurement
 - Better apparatus
 - Better procedure
 - Repeat

□Estimate







Systematic Error

- Determines the *accuracy* of any measurement
- □*May* be present in every physical measurement
 - calibration
 - uniform or controlled conditions (e.g., avoid systematic changes in temperature, light intensity, air currents, etc.)
- Identify & eliminate or reduce

24.32 24.24 24.16 24.08 24 40 20 60 80 100 0 trial

chronological data

Uncertainties in Modeling & Simulation

Aleatory Uncertainty:

- *inherent randomness* in the system. Also known as:
 - stochastic uncertainty
 - variability
 - irreducible uncertainty

Epistemic Uncertainty:

- due to *lack of perfect knowledge* about the system. Also known as:
 - Incertitude
 - system error
 - reducible uncertainty



Approximations in simulation

- From mathematical models to numerical models
 - Taylor series
 - Functional analysis

□From numerical models to computer codes

- Discretization (differentiation, integration)
- Searching algorithms (solving equations, optimization)



Mathematical models → Numerical models Approximation in Taylor Series

Truncation

$$f(x) = f(x_0) + f'(x_0)(x - x_0) + O(h^2)$$

= $f(x_0) + f'(x_0)h + O(h^2)$

$$f(x) = f(x_0) + f'(x_0)h + \frac{1}{2}f''(x_0)h^2 + O(h^3)$$



Convert complex functions into simple and computable ones by transformation in vector spaces

- Fourier analysis
- Wavelet transform
- Polynomial chaos expansion
- Spectral methods
- Mesh-free methods

• • •



□Approximate the original f(x) by linear combinations of basis functions $\psi_i(x)$'s as

$$f(x) \approx \sum_{i=0}^{N} c_{i} \psi_{i}(x)$$

□In a vector space (e.g. Hilbert space) with an infinite number of dimensions

$$f(x) = \sum_{i=0}^{\infty} c_i \psi_i(x)$$



□An *inner product* $\langle f, g \rangle$ is defined as a "projection" in the vector space, such as

$$\left\langle f,g\right\rangle \coloneqq \int_{-\infty}^{\infty} f(x)g(x)W(x)dx$$

Typically we choose orthogonal basis functions $\psi_i(x)$'s such that

$$\left\langle \psi_i, \psi_j \right\rangle = \int_{-\infty}^{\infty} \psi_i(x) \psi_j(x) W(x) dx = \begin{cases} Constant & (\forall i, j, i = j) \\ 0 & (i \neq j) \end{cases}$$

for orthonormal basis functions

$$\left\langle \psi_i, \psi_j \right\rangle = \int_{-\infty}^{\infty} \psi_i(x) \psi_j(x) W(x) dx = \begin{cases} 1 & (i=j) \\ 0 & (i\neq j) \end{cases}$$

 \Box The coefficients c_i 's are computed by

$$c_{_{i}}=\frac{\left\langle f,\psi_{_{i}}\right\rangle }{\left\langle \psi_{_{i}},\psi_{_{i}}\right\rangle }$$

□The computable function is

$$f(x) \approx \sum_{i=0}^{N} c_i \psi_i(x)$$

with *truncation*!



Numerical Model → Computer Code Compute integrals

□Quadrature

- Approximate the integrand function by a polynomial of certain degree
- Approximate the integral by the weighted sum of regularly sampled functional values





Numerical Model → Computer Code Compute integrals

Monte Carlo simulation

- Let p(u) denote uniform density function over [a, b]
- Let U_i denote i th uniform random variable generated by Monte Carlo according to the density p(u)
- Then, for "large" N

$$\int_{a}^{b} f(x) dx \approx \frac{b-a}{N} \sum_{i=1}^{N} f(U_i)$$

Variance reduction (importance sampling) to improve efficiency



Numerical Model → Computer Code Compute derivatives

□Finite-divided-difference methods

- Approximated derivatives come from Taylor series
 - e.g. forward-finite-difference

$$f(x_{i+1}) = f(x_i) + f'(x_i)(x_{i+1} - x_i) + O(h^2) = f(x_i) + f'(x_i)h + O(h^2)$$
$$f'(x_i) = \frac{f(x_{i+1}) - f(x_i)}{h} + O(h)$$



Floating-Point Representation

□How does computer represent numbers?



Perfect world



Imperfect world



Ariane 5

- Exploded 37 seconds after liftoff
- Cargo worth \$500 million

Why

- Computed horizontal velocity as floating point number
- Converted to 16-bit integer
- Worked OK for Ariane 4
- Overflowed for Ariane 5
 - Used same software





Do you trust your computer?

Rump's function: $f(x, y) = 333.75y^6 + x^2(11x^2y^2 - y^6 - 121y^4 - 2) + 5.5y^8 + \frac{x}{2y}$ f(x = 77617, y = 33096) = ?Single precision: f = 1.172603...Double precision: f = 1.1726039400531...Extended precision: f = 1.172603940053178...Correct one is: f = -0.8273960599468213











Another Story



□On February 25, 1991

- A Patriot missile battery assigned to protect a military installation at Dhahran, Saudi Arabia
- But ... failed to intercept a Scud missile28 soldiers died
- **u**... an error in computer arithmetic

 $0.1 \times 10 \neq 1$



IEEE 574 Standard

IEEE Floating Point Representation

	s	exponent	E	mantissa	T				
	1 bit	w bits	8 bits		23 bits	t = p - 1 bits			
	IEEE Double Precision Floating Point Representation								
	1 bit		11 bits		52 bits				
	s	exponent		mantissa					
□ If $E = 2^{w}$ -1 and $T \neq 0$, then v is NaN regardless of S.									
\Box If $E = 2^{w} - 1$ and $T = 0$, then $v = (-1)^{S \times \infty}$.									
$\Box \text{ If } 1 \leq E \leq 2^{w} - 2, \text{ then } v = (-1)^{S} \times 2^{E - bias} \times (1 + 2^{1 - p} \times T);$									
normalized numbers have an implicit leading significand bit of 1. □ If $E = 0$ and $T \neq 0$, $v = (-1)^{S} \times 2^{emin} \times (0 + 2^{1-p} \times T)$;									
denormalized numbers have an implicit leading significand bit of 0.									
\Box If $E = 0$ and $T = 0$, then $v = (-1)^{S \times 0}$ (signed zero)									
where $bias=2^{w-1}-1$ and $emin=2-2^{w-1}=1-bias$									

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Distribution of Values

□6-bit IEEE-like format

- w = 3 exponent bits
- t = 2 fraction/mantissa bits
- bias = 3





Distribution of Values (zoom-in view)

□6-bit IEEE-like format

- w = 3 exponent bits
- t = 2 fraction/mantissa bits

• bias = 3





Round-Off Errors



Overflow error – "not large enough"
Underflow error – "not small enough"
Rounding error – "chopping"

http://www.cs.utah.edu/~zachary/isp/applets/FP/FP.html



Rounding

$$D \times \beta^{E} = \pm \left(d_0 + \frac{d_1}{\beta} + \frac{d_2}{\beta^2} + \dots + \frac{d_{p-1}}{\beta^{p-1}} \right) \beta^{E}$$

- Round by chopping (round toward zero)
 - Truncate base expansion after (*p*-1)st digit
 - Machine epsilon $\varepsilon_{machine} = \beta^{1-p}$
- Round to nearest (round to even)
 - Last digit is even in case of tie
 - Machine epsilon $\varepsilon_{machine} = \frac{1}{2}\beta^{1-p}$ $\left|\frac{fl(x) x}{x}\right| \le \varepsilon_{machine}$



Cancellation

Subtraction between two *p*-digit numbers having the same sign and similar magnitudes yields result with fewer than *p* digits.

Significant digits of two numbers cancel.
 Despite exactness of result, cancellation often implies serious loss of information.
 Relative uncertainty in difference is largely due to previous rounding errors.

 (1+ε)-(1-ε)=1-1=0

Special Numbers

Expression	Result
0.0 / 0.0	NaN
1.0 / 0.0	Infinity
-1.0 / 0.0	-Infinity
NaN + 1.0	NaN
Infinity + 1.0	Infinity
Infinity + Infinity	Infinity
NaN > 1. 0	false
NaN == 1.0	false
NaN < 1.0	false
NaN == NaN	false
0.0 == -0.0	true

standard range of values permitted by the encoding (from 1.4e-45 to 3.4028235e+38 for float)

Floating point Hazards

This expression	does NOT equal to this expression	when
0.0 - f	-f	f is o
f < g	! (f >= g)	f or g is NaN
f == f	true	f is NaN
f + g - g	f	g is infinity or NaN

double s=0; for (int i=0; i<26; i++) s += 0.1; System.out.println(s);

double d = 29.0 * 0.01; System.out.println(d); System.out.println((int) (d * 100)); The result is0.2928



Comparing Floating Point Numbers

Try to avoid floating point comparison directly
Testing if a floating number is greater than or less than zero is even risky.

- □*Instead*, you should compare the absolute value of the difference of two floating numbers with some pre-chosen epsilon value, and test if they are "close enough"
- □If the scale of the underlying measurements is unknown, the test "abs(*a*/*b* 1) < epsilon" is more robust.

Don't use floating point numbers for exact values



Uncertainties in M&S

Model errors due to approximations in truncation or sampling

- Taylor approximation
- Functional analysis

Numerical errors due to floating-point representation

Round-off errors



Summary

 Modeling is abstraction
 M&S always has approximations involved, which are important sources of epistemic uncertainty.

Computer tricks us

Abstraction Approximation Round-off



Further Readings

□ Goldberg, D. (1991) "What every computer scientist should know about floating-point arithmetic," *ACM Computing Surveys*, **23**(1), 5-48

